## What is a sheaf?



## Michael Robinson

## DARPA SIMPLEX Program


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## Session objectives

- What is heterogeneous data fusion?
- What is a sheaf?
- What happens if I don't have a sheaf?


## Axiom 1: There is a set of data sources

- A list of sensors
- A collection of tables

Grad school

High school


Postdoc

## Axiom 2: There is a list of possible attributes

Different sensors will read out in attributes

- Possible columns or keys

- Possible signal spaces


Grad school


Postdoc

Axioms 3 and 4: The data sources are topologized

- Topology comes from shared attributes



## Axiom 5: Data sources can be compared via restriction

- Essentially by forming subtables
- Transformations are permitted



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## Axiom 5: Data sources can be compared via restriction

- Essentially by forming subtables
- Transformations are permitted



## Axioms 6 and 7: Data sources can be combined in a unique way

- Table joins: can find rows of a bigger table
- Problems can arise if this isn't satistifed!



## Axioms 6 and 7: Data sources can be combined in a unique way

- Table joins:
- Might result in empty tables (no rows)



## Axioms 6 and 7: Data sources can be combined in a unique way

- Table joins:
- Might result in empty tables (no rows)



## Mathematical representation

- Specifies the attribute spaces
- Specifies transformations


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## Key/table duality

Two obvious viewpoints:

- Table-centric $\rightarrow$ Tables on vertices (leads to sheaf)
- "Bottom up"
- Restrictions go from low dimensional simplices to higher dimensional simplices
- Key-centric $\rightarrow$ Keys on vertices (leads to cosheaf)
- "Top down"
- Extensions go from high dimensional simplices to lower dimensional simplices


## Building a cosheaf model

- Keys are vertices, and tables are simplices (so the base space is different)



## What happens when the axioms fail?

- Axiom 1 or 2: Data isn't in a set...
- You won't be able to do much computationally
- Axiom 3 or 4 : Sources not topologized
- No basis for combining data sources...
- Axiom 5: No transformations to identify commonalities between observations
- Although information about a given entity might be available through different sources, they can't be joined
- Axiom 6 or 7: Cannot uniquely fuse
- Even if two observations are comparable, it is impossible to infer anything else about nearby observations
- This often happens in databases - two rows with overlapping keys and matching values doesn't mean they're the same!


## Now, abstractly...

## A sheaf of

$\qquad$ on a $\qquad$
(data type)
(topological space)

## Simplicial complexes

- ... higher dimensional simplices (tuples of vertices)



## Simplicial complexes

- The attachment diagram shows how simplices fit together



## A sheaf is ...

- A set assigned to each simplex and ...



## Each such set is called the

 stalk over its simplex$\mathbb{R}^{2}$
This is a sheaf of vector spaces on a simplicial complex

## A sheaf is ...

- ... a function assigned to each simplex inclusion



## A sheaf is ...

- ... so the diagram commutes.



## A global section is ...

- An assignment of values from each of the stalks that is consistent with the restrictions



## Some sections are only local

- They might not be defined on all simplices or disagree with restrictions



## Flabbiness

- If all local sections defined on vertices extend to global sections, the sheaf is called flabby (or flasque)
- These sheaves don't have interesting invariants
- They are good for decomposing other sheaves

$$
\mathbb{R}^{3} \rightarrow 0 \longleftarrow \mathbb{R}^{3}
$$

- Example: Vertex- or edge-weighted graphs with no further constraints
- Flabby sheaves mean there is a lack of constraints imposed by the model


## Queue model example

## A queue as a sheaf

- Contents of the shift register at each timestep
- $N=3$ shown

$$
\left(\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0
\end{array}\right) \quad\left(\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0
\end{array}\right) \quad\left(\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0
\end{array}\right)
$$

$\rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}$

$$
\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right) \quad\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right) \quad\left(\begin{array}{ll}
0 & 1
\end{array}\right)
$$

## A single timestep

- Contents of the shift register at each timestep
- $N=3$ shown

$$
\left(\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0
\end{array}\right) \quad\left(\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0
\end{array}\right) \quad\left(\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0
\end{array}\right)
$$

$$
\begin{aligned}
& (9,2) \leftrightarrow(1,9,2) \rightarrow(1,9) \leftarrow(1,1,9) \rightarrow(1,1) \\
& \left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right) \\
& \left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right) \\
& \left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right)
\end{aligned}
$$

## Wireless network example

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## Wireless activation sheaf

- Each node in the wireless complex has a unique ID
- The stalk over a face consists of the set of nodes adjacent to that face, and the special symbol $\perp$



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## Wireless activation sheaf

- Each node in the wireless complex has a unique ID
- The stalk over a face consists of the set of nodes adjacent to that face, and the special symbol $\perp$

Note: "adjacent" means "has a higher-dimensional face in common"


## Wireless activation sheaf

- Restrictions map node IDs via identity whenever possible, and otherwise send to $\perp$
restriction



## Wireless activation sheaf

- Sheaves model node activation and traffic-passing protocols, by encoding local consistency constraints
- Node activation patterns when interference is possible. ( $\perp$ means no activity)


Link complex
Some possible sections

$$
1 \rightarrow 1 \leftarrow 1 \rightarrow \perp \leftarrow \perp
$$


$\{\perp, 1,2\} \quad\{\perp, 1,2,3\} \quad\{\perp, 2,3\}$
Activation sheaf

# Shared situational awareness 

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## Forming mosaics

- Multiple, overlapping images can be assembled into a mosaic by stitching together similar regions
- Many algorithms exist
- Most are robust to perspective (or other) changes


Overlapping maps

(Image courtesy of NASA/JPL)

# Shared visual situational awareness (In collaboration with UCLA) 




## Heterogeneous fusion among homogeneous sensors



## Heterogeneous fusion among homogeneous sensors


"Physical" sensor footprints
Sensor data space

## Heterogeneous fusion among homogeneous sensors

This construction - the data together with the transformations - is a sheaf $\longrightarrow$ Fused targets


Camera 1
Camera 2

## Data Structures as Sheaves



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## (DARPA SIMPLEX Program


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## Session objectives

- How do sheaves extend well-known data structures?
- How do I translate between sheaf-based data structures?
- What can I do once I have a sheaf?


## What is a sheaf?



## What is a sheaf?



## What isn't a sheaf?

If labels on the graph are not systematically related to one another...

Consistency checks become ad hoc


Cross-modality inference is no longer possible


## Vertex- or (hyper)edge-weighted (hyper)graphs

Vertex weighted $\rightarrow$ sheaves

- Vertex has nontrivial stalk
- All restrictions are zero maps
- The resulting sheaf is flabby

Hyperedge-weighted $\rightarrow$ cosheaves

- Toplex has nontrivial stalk
- All extensions are zero maps
- The resulting cosheaf is coflabby*


## Flow sheaves

- Start with a collection of paths along which material flows
- Label each track segment with amount of material on that segment



## Flow sheaves

- Conservation law enforced at each vertex
- Depending on precisely how we count material (in $\mathbb{N}$ or $\mathbb{R}$, for instance), we might write the conservation law as



## Flow sheaf

- Each degree $n$ vertex is assigned a free $s R$-module* of rank $(n-1)$ for material measured in a semiring $s R$
- Restriction maps are projections

* or just a vector space if it seems easier


## Local consistency

A flow sheaves encodes a notion of consistency between adjacent faces


## Inferential ambiguity

Depending on how we make measurements, we might not get "the full story" of the flow

Many possible inferences


Exactly one inference

## Bayesian networks

## Probability spaces

Start with a set of random variables $X_{0}, X_{1}, \ldots X_{n}$
Consider the set $P\left(X_{0}, X_{1}, \ldots, X_{n}\right)$ of all joint probability distributions over these random variables

- These are the nonnegative measures (generalized functions)

$$
f=f\left(X_{0}, X_{1}, \ldots, X_{n}\right)
$$

with unit integral

- This is not a vector space - adding probability distributions doesn't yield another distribution


## Marginalization cosheaf

There is a natural map

$$
P\left(X_{0}, X_{1}, \ldots, X_{n}\right) \rightarrow P\left(X_{0}, X_{1}, \ldots, X_{n-1}\right)
$$

via marginalization, namely

$$
f\left(X_{0}, X_{1}, \ldots, X_{n-1}\right)=\int f\left(X_{0}, X_{1}, \ldots, X_{n}\right) d X_{n}
$$

- There similar maps for marginalizing out the other random variables, too
- This yields a cosheaf on the complete $n$-simplex!


## Bayes' rule

Conditional probabilities produce maps going the other way... For instance,

$$
P\left(X_{0}, X_{1}, \ldots, X_{n-1}\right) \rightarrow P\left(X_{0}, X_{1}, \ldots, X_{n}\right)
$$

is parameterized by functions $C$

$$
F\left(X_{0}, X_{1}, \ldots, X_{n}\right)=C\left(X_{0}, X_{1}, \ldots, X_{n}\right) f\left(X_{0}, X_{1}, \ldots, X_{n-1}\right)
$$

usually, one writes the arguments to $C$ like

$$
C=C\left(X_{n} \mid X_{0}, X_{1}, \ldots, X_{n-1}\right)
$$

So... conditional probabilities yield a sheaf on part of the $n$-simplex

## Small Bayes net example

- Consider two binary random variables $X$ and $Y$ with a given conditional $C(Y \mid X)$

$$
P(X) \stackrel{\left(\begin{array}{llll}
1 & 1 & 0 & 0 \\
0 & 0 & 1 & 1
\end{array}\right)}{\longleftrightarrow} P(X, Y) \xrightarrow{\left(\begin{array}{llll}
1 & 0 & 1 & 0 \\
0 & 1 & 0 & 1
\end{array}\right)} P(Y)
$$

Marginalization cosheaf
$P(X) \xrightarrow[\left(\begin{array}{cc}p(0 \mid 0) & 0 \\ p(110) & 0 \\ 0 & p(0 \mid 1) \\ 0 & p(111)\end{array}\right)]{ }=C(X, Y)$
Conditional sheaf

# Linear translation-invariant filters 

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## How does a sheaf model a signal?



## How does a sheaf model a signal?



## A sheaf morphism ...

- ... takes data in the stalks of two sheaves ...



## A sheaf morphism ...

- ... and relates them through linear maps ...



## A sheaf morphism ...

- ... so the diagram commutes!



## A sampling morphism is ...



## An ambiguity sheaf is ...

- The collection of kernels of these sampling maps



## An ambiguity sheaf is ...



## The general sampling theorem

- Given a sheaf $S$ and a sampling of it, construct the ambiguity sheaf $A$

Theorem:

- Perfect reconstruction of global sections of $S$ is possible if and only if the only global section of $A$ is the zero function
- "No ambiguities means it's possible to reconstruct"


## Nyquist-Shannon sampling

- Encode signals as a sheaf of bandlimited functions $B F$ over $\mathbb{R}$, with bandwidth $B$
- It's easier to work in the frequency domain:

$$
B F=\{f \in \mathrm{C}(\mathbb{R}, \mathbb{C}) \mid \operatorname{supp} f \subseteq[-B, B]\}
$$

- Samples are taken at integers, obtained by inverse Fourier transform
- For instance at $n$, we sample using the function $M_{n}$

$$
M_{n}(f)=\int_{-B}^{B} f e^{-2 n \pi i x} d x
$$

## Nyquist-Shannon sampling

- The ambiguity sheaf identifies bandlimited functions with zeros at specific locations

$$
\begin{aligned}
A_{n} & =\left\{f \in B F \backslash M_{n}(f)=0\right\} \\
& =\{\text { Bandlimited functions that are zero at } n\}
\end{aligned}
$$



## Nyquist-Shannon sampling

- The ambiguity sheaf identifies bandlimited functions with zeros at specific locations

$$
A_{n}=\left\{f \in B F \mid M_{n}(f)=0\right\}
$$

$=\{$ Bandlimited functions that are zero at $n\}$


Global sections of the ambiguity sheaf are bandlimited functions that vanish at every integer


## Filters as sheaf morphisms

- Theorem: Every discrete-time LTI filter can be encoded as a sequence of two sheaf morphisms


Sheaf formalism
Input —— Internal state - Output
Hardware


## Input sheaf

- Sections of this sheaf are timeseries, instead of continuous functions

$$
\longrightarrow 0 \longleftarrow \mathbb{R} \longrightarrow 0 \longleftarrow \mathbb{R} \longrightarrow 0 \longleftarrow \mathbb{R} \longrightarrow 0
$$

## Output sheaf

- The output sheaf is the same



## The internal state

- Contents of the shift register at each timestep
- $N=3$ shown

$$
\left(\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0
\end{array}\right) \quad\left(\begin{array}{lll}
1 & 0 & 0 \\
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$$

$\rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}$

$$
\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right) \quad\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right) \quad\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right) \quad\left(\begin{array}{lll}
0 & 1 & 0 \\
0 & 0 & 1
\end{array}\right)
$$

## The internal state

- Loads a new value with each timestep




## The internal state

- Produces average of the shift register at each timestep


A


## Finishing both morphisms

- Put in a few zero maps!

$\rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}^{3} \rightarrow \mathbb{R}^{2} \leftarrow \mathbb{R}$

$0 \longleftarrow \mathbb{R} \longrightarrow 0 \longleftarrow \mathbb{R} \longrightarrow 0 \longleftarrow \mathbb{R} \longrightarrow 0 \longleftarrow \mathbb{R}$


## A single timestep

- Sections of the sheaves linked together give possible input/output combinations of the filter



## Of course, this extends...

- Sections of the sheaves linked together give possible input/output combinations of the filter

$\rightarrow(9,2) \leftarrow(1,9,2) \rightarrow(1,9) \leftarrow(1,1,9) \rightarrow(1,1) \leftarrow(5,1,1) \rightarrow(5,1) \leftarrow(2,5$



$0 \longleftarrow 4 \longrightarrow 0 \longleftarrow 3.7 \longrightarrow 0 \longleftarrow 2.3 \longrightarrow 0 \longleftarrow 2$.


## How this formalism helps...

- Of course, it corresponds nicely to the hardware ... BUT...
- It's easy to splice in nonlinear operations
- It works on nontrivial base spaces: A systematic study of filtering is now possible on cell complexes

Sheaves make it easy to invent topological filters that have controlled performance characteristics

## Filtering sinosoids from noise

Signals at ports of standard LPF filter


## Filtering out chirpy signals

Signals at ports of standard LPF filter


Time

## Filtering out chirpy signals

Signals at ports of variable-bandwidth LPF filter


Time

## Filtering out chirpy signals

Signals at ports of variable-bandwidth LPF filter


## Circumventing bandwidth limits

- More averaging in a connected window leads to:
- More noise cancellation (Good)
- More distortion to the signal (Bad)

- Can safely do more averaging by collecting samples at "similar places" across the entire signal



## Filter block diagram



# Stage 1: Topological estimation 



# Stage 2: Grouping sheaf 

## Distance matrix of point cloud



Values of the signal at the neighbors


Process execution

## Topological filter results

Signals at ports of topological filter


## Compare: standard adaptive filter

Signals at ports of variable-bandwidth LPF filter


Time

## Input image



Noisy input image


## Fixed frequency image filter

Boxcar filtered output


# Topological filter output 

Topological filtered output


## Error contributions

Topological filter stage-wise performance at $2.5 x$ filter BW


## Error contributions

Topological filter stage-wise performance at $2.5 x$ filter $B W$


## Error contributions

Topological filter stage-wise performance at $2.5 x$ filter $B W$


## Further reading...

- Sanjeevi Krishnan, "Flow-cut dualities for sheaves on graphs," http://arxiv.org/abs/1409.6712
- Robert Ghrist and Sanjeevi Krishnan, "A Topological Max-Flow Min-Cut Theorem," Proceedings of Global Signals. Inf., (2013).
- Michael Robinson, "Understanding networks and their behaviors using sheaf theory," IEEE Global Conference on Signal and Information Processing (GlobalSIP) 2013, Austin, Texas.
- Michael Robinson, "The Nyquist theorem for cellular sheaves," Sampling Theory and Applications 2013, Bremen, Germany.

