Thm 2.8: An acyclic digraph, $D$ has a unique vertex basis consisting of all vertices with no incoming arcs.

Proof: Let $B$ be the set of all vertices with no incoming arcs. If $v \in B$. Then since $v$ has no incoming arcs, $v$ is only reachable from itself. Thus $B$ must be a subset of any vertex base. Suppose $u \notin B$. Let $u_{1}=u$. Then $u_{1}$ has an incoming $\operatorname{arc}\left(u_{2}, u_{1}\right)$. If $u_{2} \in B$, then $u_{1}$ is reachable from a vertex in $B$. If $u_{2} \notin B$ then $u_{2}$ has an incoming $\operatorname{arc}\left(u_{3}, u_{2}\right)$.

Suppose the path $u_{n}, \ldots, u_{1}$ is defined such that all vertices are distinct.

If $u_{n} \in B$, then $u$ is reachable from a vertex in $B$.
If $u_{n} \notin B$ then $u_{n}$ has an incoming arc $\left(u_{n+1}, u_{n}\right)$. If $u_{n+1}=u_{i}$ for some $i=1, \ldots, n$, then $u_{n+1}, u_{n}, \ldots, u_{i}$ is a cycle, a contradiction. Hence all the vertices of $u_{n+1}, u_{n}, \ldots, u_{1}$ are distinct.

Since the number of vertices of $D$ is finite, this process must eventually end, say with the path $u_{t}, \ldots, u_{1}$. Since we cannot continue this process, $u_{t}$ must not have any incoming arcs. Hence $u_{t} \in B$, and hence $u$ is reachable from a vertex in $B$. Thus any vertex basis must be contained in $B$. Hence $B$ is the unique vertex basis of D.

